# ANOTHER METHOD OF LOGIC SYNTHESIS OF DIGITAL COUNTING CIRCUITS

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**Abstract.** In order to synthesize automat (in this case digital counters), the minimizing internal states is of particular significance and plays a decisive role in the results of synthetic circuit. This can be done in many ways, but the use of Karnaugh map is considered optimal. However, this process has some disadvantages that it can not be overcome when the number of input variants is large. In experience, if the number of variants is 7, manual minimization of circuit functions using Karnaugh map arises many difficulties and even become impossible if over 10 variants are available.

In order to deal with this weakness, it is both necessary and rational to use computer in logical synthesis of counting circuit. This is the aim of this article.

## 1. Synthesizing counting circuits using similar forms

For the method of synthesizing digital counting circuits using computers is firstly primarily based on the results achieved through the Karnaugh map [1]. By there results drawning the general laws of circuit functions for each synchronous and asynchronous counters, for each Flip-Flop (FF) and for each kind of codes. *These general laws* help to develop *mathematical models* as well as *computer programmes* which enable the fastest definition of minimized circuit functions of each desired counters.

Let us investigate, for example, the input states  $R_2$ ,  $S_2$  and outputs states  $Q_2$  of RS – FF in synchronous counters, real binary, 4 inputs (k = 4). The input states  $R_2$ ,  $S_2$  as well as outputs states  $Q_2$  are given in table 1.

**Table 1.** The input states  $R_2$ ,  $S_2$ .  $\epsilon$  is counting state,  $\epsilon = 2^k - 1 = 2^4$ - 1 = 15 = m-1, with m is a cardinal number  $Q_2$  is an output state corresponding inputs states  $S_2$  and  $R_2$ 

ε	$\mathbf{Q}_2$	$\mathbf{R}_2$	$\mathbf{S}_2$	ε	$\mathbf{Q}_2$	$\mathbf{R}_2$	$\mathbf{S}_2$	ε	$\mathbf{Q}_2$	$\mathbf{R}_2$	$\mathbf{S}_2$	ε	$\mathbf{Q}_2$	$\mathbf{R}_2$	$\mathbf{S}_2$
0	0	d	0	4	0	d	0	8	0	d	0	12	0	d	0
1	0	0	1	5	0	0	1	9	0	0	1	13	0	0	1
2	1	0	d	6	1	0	d	10	1	0	d	14	1	0	d
3	1	1	0	7	1	1	0	11	1	1	0	15	1	1	0

From table 1, we can form impulse diagrams for both  $Q_2$ ,  $R_2$  and  $S_2$  (figure 1)



Figure 1: Impulse diagram of Q<sub>2</sub>, R<sub>2</sub>, and S<sub>2</sub>. The dots represent non-defined state 'don't care', which receive either value 1 or 0 when circuit function is reduced and which can be used or not be depending on certain cases

In this case we particularly study the input state  $R_2$ .

If in minimizing circuit functions using Karnaugh map state 'don't care' is given value 0, we can have an impulse diagram of  $R_2$  as in figure 2.



Figure 2: Impulse diagram of  $Q_2$  and  $R_2$  to describe forms of circuit functions of  $R_2$ 

From figure 2 we can see that if  $Q_2$  and  $R_2$  impulses are the same state, the responding circuit functions have the same form, here called form 1. If  $Q_2$  impulse flank positive but  $R_2$  impulse flank negative, then circuit functions have the same form 2. Figure 3 describes the appearance of form 1 and form 2 on  $\varepsilon$  axis.



Figure 3: Description of form 1 and form 2 on  $\varepsilon$  axis

From figure 3 we can see that if  $\varepsilon = 2$ , 6, 10, 14..., the corresponding impulse diagrams have the same for (form 1) and if  $\varepsilon = (3, 4, 5)$ , (7, 8, 9)... the corresponding impulse diagrams have the same form (form 2). Therefore we say impulse diagrams corresponding  $\varepsilon$  in a definite period are similar.

The concept of 'similar' is understood in the following way: Two similar circuit functions are two functions that have the same mathematical structure, such as the same "sum of products" or "product of sums", or similar sum, similar product, but the connotations are different.

This concept of 'similar' is the basis on which mathematical models for circuit functions of counters, in this case synchronous RS-FF counter, real binary code, are formed.

Also from table 1, using Karnaugh map method, we can define circuit functions corresponding to inputs  $R_2$  (Table 2).

Table 2: circuit functions corresponding to input  $R_2$  of RS - FF - Counter, k = 4,  $\epsilon = 0$  to 15

3	$ m R_2$	3	$R_2$	3	$R_2$	3	$ m R_2$
2	$\mathbf{Q}_2$	6	$Q_1Q_2$ + $Q_2Q_3$	10	$Q_1Q_2$ + $Q_2Q_4$	14	$\mathbf{Q}_1\mathbf{Q}_2$ + $\mathbf{Q}_2\mathbf{Q}_3\mathbf{Q}_4$
3	$\mathbf{Q}_1\mathbf{Q}_2$	7	$\mathbf{Q}_1\mathbf{Q}_2$	11	$\mathbf{Q}_1\mathbf{Q}_2$	15	$\mathbf{Q}_1\mathbf{Q}_2$
4	$\mathbf{Q}_1\mathbf{Q}_2$	8	$\mathbf{Q}_1\mathbf{Q}_2$	12	$\mathbf{Q}_1\mathbf{Q}_2$		
5	$\mathbf{Q}_1\mathbf{Q}_2$	9	$\mathbf{Q}_1\mathbf{Q}_2$	13	$\mathbf{Q}_1\mathbf{Q}_2$		

From figure 2, figure 3, table 2 and the concept of "similar" we can see that circuit functions have two main forms:

Form 1:

$$\mathbf{R}_{l,\varepsilon}^{(1)} = \mathbf{Q}_1 . \mathbf{Q}_2 ... . \mathbf{Q}_l = \prod_{i=1}^{\ell} \mathbf{Q}_i = \mathbf{A} \qquad \forall \varepsilon \in \{2, 6, 10, 14\}$$
(1)

Form 2:

$$R_{1,\varepsilon}^{(2)} = A + \prod_{i=1}^{t} Q_i = A + B, \quad \forall \varepsilon \notin \{2, 6, 10, 14\}$$
(2)

with

$$\mathbf{B} = \prod_{i=1}^{t} \mathbf{Q}_i \tag{3}$$

If we can prove that the existence of form 1 and 2 follows a certain law, for example form 1 and 2 exist at the same time in a repetitive period with denominator  $\Delta \varepsilon = 2^{\ell} = 4$  (in this case  $\ell = 2$ ):  $\varepsilon_1 = \{(3, 4, 5), (7, 8, 9), (11, 12, 13)...\};$  $\varepsilon_2 = \{2, 6, 10, 14...\}$  and if we can *define the circuit functions* of form 1 (as well as form 2) with  $\varepsilon_1 = (3, 4, 5)$  (as well as  $\varepsilon_2 = 2$ ), we can define all circuit functions of input  $R_2$  (as well as  $S_2$ ) with any  $\varepsilon$  counting state.

## 1.1 Prove the existence of circuit functions follows a certain law

The problem is we have to prove with any unlimited great k variable input, that is with any unlimited great  $\varepsilon$  counting state, the impulse diagram always follows a certain law, that is always exist form 1 and 2 corresponding a repetitive period with  $\Delta \varepsilon = 2\ell$  (in the case of counter RS=FF).

Assuming that  $f(\mathbf{\epsilon}) = \mathbb{R}^{(1)}_{1,\mathbf{\epsilon}}$  is a function satisfying term Dirichlet of Fourier theorem on period [3, 4, 5] = [a,b]. In order to develop  $f(\mathbf{\epsilon})$  into Fourier series, we form a periodic function  $g(\mathbf{\epsilon})$  having a period either bigger or equal to (b - a) so that  $g(\mathbf{\epsilon}) = f(\mathbf{\epsilon}), \forall \mathbf{\epsilon} \in [a,b]$ 

Obviously there are many ways to develop  $g(\boldsymbol{\epsilon})$  into Fourier series. For each  $g(\boldsymbol{\epsilon})$  there is a corresponding Fourier series, therefore there are a number of Fourier series demonstrating  $f(\boldsymbol{\epsilon}) = R_{1,\epsilon}^{(1)}$ . Similarly,  $f(\boldsymbol{\epsilon}) = R_{1,\epsilon}^{(2)}$  can also be developed into a Fourier series. To put it simple, the circuit function  $f(\boldsymbol{\epsilon}) = R_{\ell,\epsilon}^{(1)} + R_{\ell,\epsilon}^{(2)}$  with every  $\boldsymbol{\epsilon}$  is periodicall with period  $\Delta \boldsymbol{\epsilon} = 2^{\ell}$  (in this case  $\ell = 2 \rightarrow \Delta \boldsymbol{\epsilon} = 4$ ) (Figure 4).



**Figure 4:** Circuit function developed into Fourier series with  $\Delta \varepsilon = 4$ 

Now that we can assert that with any variable input  $\ell$ , that is the counter can (theoretically) count to infinite number, then the impulse diagram of circuit function change periodically in those periods which have similar impulses, that is circuit functions always have form 1 and 2 according to certain  $\Delta \epsilon$ .

We particularly study the characters of this law for counters, firstly the circuit function in form 2.

In order to identify and analyze the forming of form 2 of input state  $R_{\ell,\epsilon_{,}}$  we investigate the circuit function of for example  $R_3$  in RS - FF - Counter with k = 6 (table 3).

**Table 3.** Circuit function with input state  $R_{3,\epsilon}$  of RS - FF - Counter with k = 6,  $\epsilon = 0$  to 64, in which: P is the periodical existence of circuit function form 2; F is the frequency of existence of circuit function in each period P, with corresponding denominator  $\Delta \epsilon = 2^{\ell}$ 

Р	F	ε	Circuit function in	form 2 of $R_3$	$\mathcal{E}_3$
	0	4	$Q_3$		$2^2 + 0.2^3 + 0$
0	1	5	$\mathbf{Q}_1 \; \mathbf{Q}_3$		$2^2 + 0.2^3 + 1$
	2	6	$\mathbf{Q}_2 \; \mathbf{Q}_3$		$2^2 + 0.2^3 + 2$
1	0	12	$Q_1 Q_2 Q_3 + Q_3 Q_4$	1	$2^2 + 1.2^3 + 0$
	1	13	$\mathbf{Q}_1  \mathbf{Q}_2  \mathbf{Q}_3 + \mathbf{Q}_1 \mathbf{Q}_3$	$_{3}\mathbf{Q}_{4}$	$2^2 + 1.2^3 + 1$
	2	14	$\mathbf{Q}_1 \ \mathbf{Q}_2 \ \mathbf{Q}_3 + \mathbf{Q}_2 \mathbf{Q}_3$	$_{3}\mathbf{Q}_{4}$	$2^2 + 1.2^3 + 2$
2	0	20	$\mathbf{Q}_1  \mathbf{Q}_2  \mathbf{Q}_3 + \mathbf{Q}_3 \mathbf{Q}_5$	5	$2^2 + 2.2^3 + 0$
	1	21	$\mathbf{Q}_1  \mathbf{Q}_2  \mathbf{Q}_3 + \mathbf{Q}_1 \mathbf{Q}_3$	$_{3}\mathbf{Q}_{5}$	$2^2 + 2.2^3 + 1$
	2	22	$\mathbf{Q}_1 \ \mathbf{Q}_2 \ \mathbf{Q}_3 + \mathbf{Q}_2 \mathbf{Q}_3$	$_{3}\mathbf{Q}_{5}$	$2^2 + 2.2^3 + 2$
3	0	28	$\mathbf{Q}_1  \mathbf{Q}_2  \mathbf{Q}_3 + \mathbf{Q}_3  \mathbf{Q}_3$	$_4\mathrm{Q}_5$	$2^2 + 3.2^3 + 0$
	1	29	$\mathbf{Q}_1  \mathbf{Q}_2  \mathbf{Q}_3 + \mathbf{Q}_1 \mathbf{Q}_3$	$_3$ $\mathbf{Q}_4 \mathbf{Q}_5$	$2^2 + 3.2^3 + 1$
	2	30	$\mathbf{Q}_1 \ \mathbf{Q}_2 \ \mathbf{Q}_3 + \mathbf{Q}_2 \mathbf{Q}_3$	$_3$ $\mathbf{Q}_4 \mathbf{Q}_5$	$2^2 + 3.2^3 + 2$
4	0	36	$\mathbf{Q}_1  \mathbf{Q}_2  \mathbf{Q}_3 + \mathbf{Q}_3 \mathbf{Q}_6$	3	$2^2 + 4.2^3 + 0$
	1	37	$\mathbf{Q}_1 \ \mathbf{Q}_2 \ \mathbf{Q}_3 + \mathbf{Q}_1 \mathbf{Q}_3$	$_{3}\mathbf{Q}_{6}$	$2^2 + 4.2^3 + 1$
	2	38	$\mathbf{Q}_1  \mathbf{Q}_2  \mathbf{Q}_3 + \mathbf{Q}_2 \mathbf{Q}_3$	$_{3}\mathbf{Q}_{6}$	$2^2 + 4.2^3 + 2$
5	0	44	$\mathbf{Q}_1 \ \mathbf{Q}_2 \ \mathbf{Q}_3 + \mathbf{Q}_3 \ \mathbf{Q}_3$	$_4\mathbf{Q}_6$	$2^2 + 5.2^3 + 0$
	1	45	$\mathbf{Q}_1 \ \mathbf{Q}_2 \ \mathbf{Q}_3 + \mathbf{Q}_1 \mathbf{Q}_3$	$_{3}\mathbf{Q}_{4}\mathbf{Q}_{6}$	$2^2 + 5.2^3 + 1$
	2	46	$Q_1 Q_2 Q_3 + Q_2 Q_3$	$_{3}$ $\mathbf{Q}_{4}\mathbf{Q}_{6}$	$2^2 + 5.2^3 + 2$
•			•		
•			•		
•					

From table 3 we can form the relation between P, F and  $\boldsymbol{\varepsilon}$  for two forms of circuit functions. (Figure 5)





**Figure 5:** Description of the relation between P, F of circuit function  $R_3$  and  $R_4$  on  $\epsilon$  axis

Circuit functions of form 1 and 2 depend on  $\varepsilon$ ,  $\ell$  ( $\ell \in k$ ). This dependence is demonstrated by the periodical existence of P and frequency of existence in each period F (See table 3 and Figure 4)

From table 3 and figure 5 we can identify the relation between  $\ell$ ,  $\epsilon$ , P and F, that is the relation as well as the mathematical models showing the existence law of circuit functions:

$$\boldsymbol{\varepsilon}_{\ell} = 2^{\ell \cdot 1} + \mathbf{P} \cdot 2^{\ell} + \mathbf{F} \tag{4}$$

From (4) we have:

$$\frac{\varepsilon_{1}}{2^{l}} = P + \frac{2^{l} + F}{2^{l}}$$
(5)

Put

$$I = I = 2^{1 - 1} + F$$
 (6)

Apply to (6) we have

$$P = \frac{\varepsilon - I}{2^{1}}$$
(7)

From (5), (6), and (7) we see that parameter I can show the complete frequency and periodical existence of circuit function.

Call E set of I from  $I_0$  to  $I_{max}$ , we have:

$$\mathbf{E} = \{\mathbf{I}_0, \, \mathbf{I}_1, \, \mathbf{I}_2, \, \dots, \, \mathbf{I}_{\max}\}$$
(8)

$$\mathbf{E} = \{\mathbf{F}_0 + 2^{\ell_1}, \, \mathbf{F}_1 + 2^{\ell_1}, \, \dots, \, \mathbf{F}_{\max} + 2^{\ell_1}\}$$
(9)

or

$$\mathbf{E} = \{ 2\ell - 1, 1 + 2\ell - 1, 2 + 2\ell - 1 \dots, 2\ell - 2 \}$$
(10)

with  $F_{max} = 2^{\ell} - 2$ 

Obviously now the value of set E shows the complete parameters of periodical as well as frequency existence of forms of circuit functions (from  $F_0$  to  $F_{max}$ ), in other words, in order to identify P and F we only have to identify the value of set E. Figure 6 shows the existence of circuit function with E corresponding  $R_3$  ( $\ell = 3$ ).



Figure 6: Describes the existence of period P and frequency F through set E

If we look into figure 6, we can see that for each  $R_3$ , if  $E \in \{4,5,6\}$  the circuit function will have form 2 and  $E \notin \{4,5,6\}$  form 1.

## 1.2 Identifying circuit function

Another problem is that we have to identify the circuit functions of input  $\ell$  for any counting state  $\varepsilon$ , in other words, we have to identify to relation between input state  $R_{\ell}$  and output state  $Q_{\ell}$  through counting state  $\varepsilon$ .

Identify circuit function form 1:

With  $A = \prod_{i=1}^{r} Q_i$ , we can obviously identify circuit function for any  $\ell \in k$ . for example the third input of RS-FF - counter, that is  $\ell = 3$ , has circuit equation as following:

$$\mathbf{R}_{3, \epsilon} = \prod_{i=1}^{3} \mathbf{Q}_{i} = \mathbf{Q}_{1} \cdot \mathbf{Q}_{2} \cdot \mathbf{Q}_{3} \qquad \boldsymbol{\epsilon} \text{ satisfies } \boldsymbol{E} \notin \{4, 5, 6\}$$

with any *l* we have:

$$\mathbf{R}_{1,\varepsilon} = \prod_{i=1}^{1} \mathbf{Q}_{i} = \mathbf{Q}_{1} \cdot \mathbf{Q}_{2} \dots \mathbf{Q}_{1} \qquad \boldsymbol{\varepsilon} \text{ satisfies } \boldsymbol{E} \notin \{2^{\ell \cdot 1}, \ 1 + 2^{\ell \cdot 1}, \ 2 + 2^{\ell \cdot 1} \dots, \ 2^{\ell} - 2\}$$

\* Identifying circuit function form 2:

The problem is to identify circuit functions  $B = \prod_{i=1}^{k} Q_i$ 

Through investigating circuit functions forms B for input R (as well as S) of counter RS – FF we have an interesting remark: for each different  $\varepsilon$  counting state there is very different circuit function, apparently not following any law (see table 3) but if we look at the relation between binary number and decimal number we will see that the binary number showing the total weight  $\eta_B$  of  $Q_\ell$  equals to the decimal number showing counting state  $\varepsilon$ , that is,  $\eta_B = \varepsilon$ .

Assuming that the counter has k inputs, then the corresponding weight to each input will be:

$\mathbf{Q}_{\mathbf{k}}$	 $Q_\ell$	 $\mathbf{Q}_5$	$\mathbf{Q}_4$	$\mathbf{Q}_3$	$\mathbf{Q}_2$	$\mathbf{Q}_1$
$\eta_{\rm k}$	 $\eta_{_\ell}$	 $\eta_{5}$	$\eta_4$	$\eta_3$	$\eta_2$	$\eta_1$
$2^{k \cdot 1}$	 $2^{\ell-1}$	 $2^4$	$2^3$	$2^2$	$2^1$	$2^0$

Now the total weight of  $Q_1$  to  $Q_k$  is

$$\eta_B = \mathbf{b}_{k+1} \cdot 2^k + \mathbf{b}_k \cdot 2^{k-1} + \dots + \mathbf{b}_2 + \mathbf{b}_1^0 = \sum_{i=1}^k \mathbf{b}_i \cdot 2^{i-1}$$
(11)

And from the conclusion  $\varepsilon = 19$ , k = 5 we have

$$\begin{split} \eta_B &= 19 = \mathbf{b}_6 \,.\, 2^5 + \mathbf{b}_5 \,.\, 2^4 + \mathbf{b}_4 \,.\, 2^3 + \mathbf{b}_3 \,.\, 2^2 + \mathbf{b}_2 \,.\, 2^1 + \mathbf{b}_1 \,.\, 2^0 \\ &= \mathbf{b}_6 \,.\, 32 + \mathbf{b}_5 \,.\, 16 + \mathbf{b}_4 \,.\, 8 + \mathbf{b}_3 \,.\, 4 + \mathbf{b}_2 \,.\, 2 + \mathbf{b}_1 \,.\, 1 \\ &= 0.32 + 1.16 + 0.8 + 0.4 + 1.2 + 1.1 \end{split}$$

That factors  $b_6$ ,  $b_4$ ,  $b_3$  must equal to 0 and  $b_5 = b_2 = b_1 = 1$  when  $\eta_B = 19$ 

In other words, when k = 5 with  $\varepsilon = 19$  the circuit function B will have  $Q_{\ell} = (Q_4, Q_1, Q_0)$ , that is,  $B_{19,5} = \prod_{i=1}^{5} Q_i = Q_0 \cdot Q_1 \cdot Q_4$ 

Now we can form a formula to find B:

$$B = \prod_{i=1}^{k} (Q_i + \overline{b_i}), \text{ with } \overline{b_i} = \begin{cases} 1 & b_i = 0\\ 0 & b_i = 1 \end{cases}$$
(12)

From this we can form the formula to identify the circuit function for each corresponding  $R_{l,\varepsilon}$  of RS - FF - Counter

$$R_{\ell,\mathcal{E}} = \Delta_{A} \cdot \prod_{i=1}^{\ell} Q_{i} + \Delta_{B} \cdot Q_{\ell} \prod_{i=1}^{t} (Q_{i} + \overline{b_{i}}) \cdot \overline{\Delta}\ell$$
(13)

In which:

$$\Delta_{\mathbf{A}} = \begin{cases} 1 & \epsilon \ge 2^{\ell} - 1 \\ 0 & \text{other cases} \end{cases}$$

 $\Delta_A$  shows the condition for the existence of form A

$$\Delta_{\mathbf{B}} = \begin{cases} 1 & \mathbf{I} = \mathbf{F} + 2^{1-1} \in \mathbf{E} = \left\{ 2^{1-1}, 1 + 2^{1-1}, \dots, 2^{1} - 2 \right\} \\ & \eta_{\mathbf{B}} = \varepsilon - 2^{\ell - 1} = \sum_{i=1}^{t} \mathbf{b}_{i} \cdot 2^{i-1} \quad \mathbf{t} = \left[ \mathrm{Id} \eta_{\mathbf{B}} + 1 \right] \\ 0 & \text{other cases} \end{cases}$$

 $\Delta_{B} \, shows the condition for the existence of form B$ 

$$\Delta_{\ell} = \begin{cases} 1 & \ell = 1 \\ 0 & \text{other cases} \end{cases}$$

 $\Delta_\ell$  shows the condition for the existence of input  $\,\ell=1$ 

The identification of  $S_{l,\varepsilon}$  is done similarly.

# 2. Conclusion

From the above results we can form a mathematical model for RS - FF -Counter and through their relation [1] (figure 7) we can identify the circuit fuction, that is the mathematical model for all counters namely JK-FF, T-FF, D-FF...



Figure 7: The relation between RS-FF with T-, JK-, D-FF

For example define the circuit function of D - FF. From figure 7 we can see that circuit function of D - FF is:

D = Sor  $D = \overline{R}$ 

That is:

$$D_{\ell,\eta} = \overline{R}_{\ell,\varepsilon} = \Delta_{A} \cdot \prod_{i=1}^{\ell} Q_{i} + \Delta_{B} \cdot Q_{\hbar} \prod_{i=1}^{t} (Q_{i} + \overline{b}_{i}) \cdot \overline{\nabla}_{\ell}$$
(14)

With these mathematical models and with software such as Pascal,  $C^{++}$ ..., and especially Mathlab we can easily design synchronous or asynchronous counters, for all codes, with any cardinal numbers using computers.

## References

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